

# Modern Userland Exploitation

From `gets()` to pwned.

# Preamble

The vulnerabilities are not real world in any way. The focus is on specific exploitation techniques. Think of it as an overly simplified model so that we may focus on other things. With this in mind, the information leak is given to us for free and an obvious stack buffer overflow is present.

# Overview

- The vulnerability
- The mitigations
- The plan
- The exploit
- Lessons
- Questions

# The Vulnerability

...

```
    gets(buf);
```

...

# The Vulnerability

```
$ man 3 gets
```

```
...
```

```
gets() reads a line from stdin into  
the buffer pointed to by s until  
either a terminating newline or EOF,  
which it replaces with a null byte  
('\0').
```

```
...
```

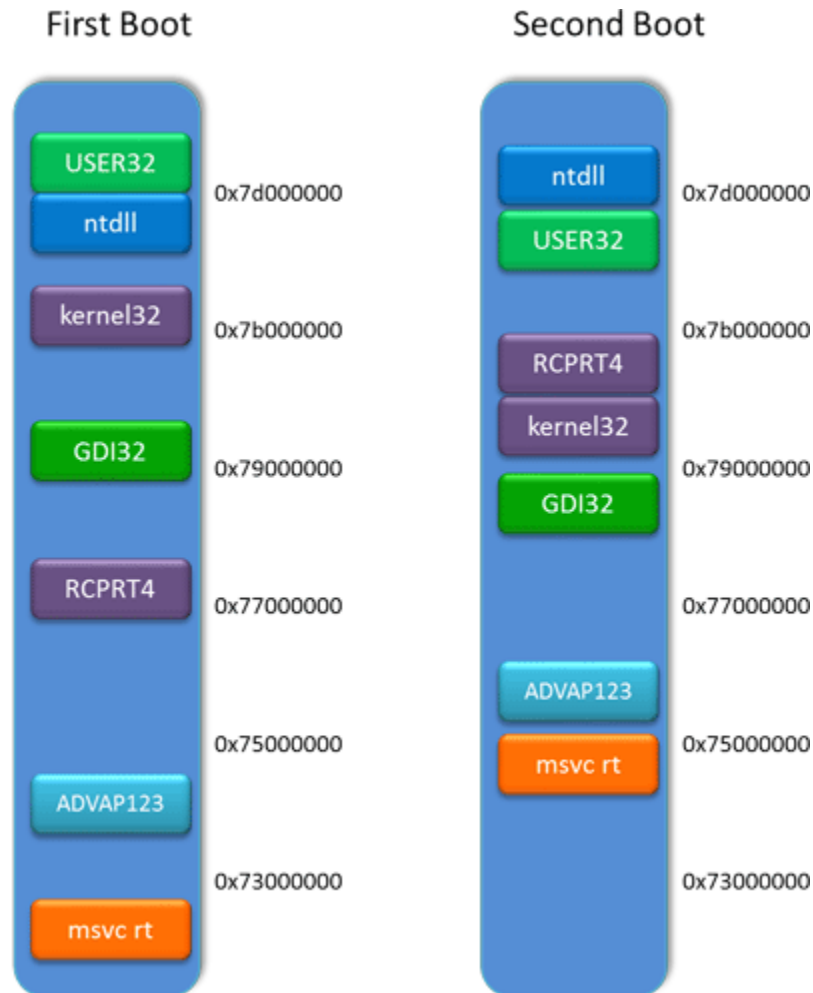
# The Mitigations

- ASLR - randomizes the base address of loaded libraries, heap, stack, and executable

# The Mitigations

Image stolen from:

<http://wasntnate.com/2012/05/how-address-space-layout-randomization-aslr-was-beaten/>



# The Mitigations

- DEP/NX/W^X - Pages in memory can have write permissions or execute permissions



# The Mitigations

## Getting around ASLR

- Information disclosure
- Using non-randomized structures
- bruteforce

## Getting around DEP

- full ROP payload
- modify protections with `mprotect()`
- create a new region with `mmap()`

# The Plan

- find the base address of a large library
- overwrite %rip
- return to mprotect() to make the stack rwx
- jmp to our, now executable, shellcode

# The Exploit

## Ingredients:

- info disclosure (bypassing ASLR)
- rop chain (bypassing DEP)
- shellcode (bypassing not executing attacker code)

# The Exploit (ASLR)

Info Disclosure (for free)

...

```
fp = fopen("/tmp/<pid>", "w");
```

...

```
fprintf(fp, "%p", dlsym(libc, "send"));
```

...

# The Exploit (DEP)

Getting around DEP with gadgets

```
$ rp -r 4 -f /lib/libc.so.6
```

```
Trying to open '/lib/libc.so.6'..
```

```
Loading ELF information..
```

```
FileFormat: Elf, Arch: Ia64
```

```
Using the Nasm syntax..
```

Wait a few seconds, rp++ is looking for gadgets..

```
in PHDR
```

```
0 found.
```

```
in LOAD
```

```
107156 found.
```

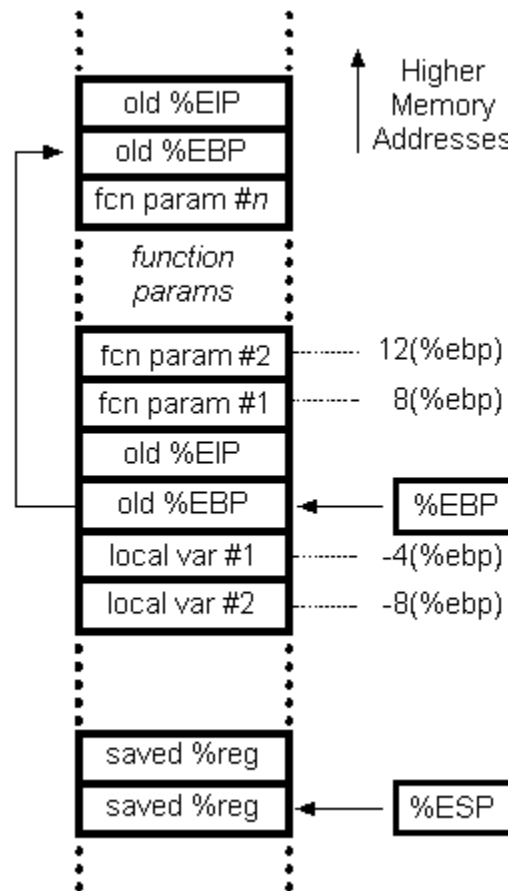
```
...
```

# The Exploit (DEP)

Before we get into using our gadgets, let's take a step back...

Here's a 32-bit call-stack:

# The Exploit (DEP)



# The Exploit (DEP)

In an exploit:

- We overwrite old %eip
- The function "returns" to old %eip (which we specified)
- We have execution control



# The Exploit (DEP)

What if we can't return to shellcode?

# The Exploit (DEP)

Why not just use a library (i.e., ret2libc)?

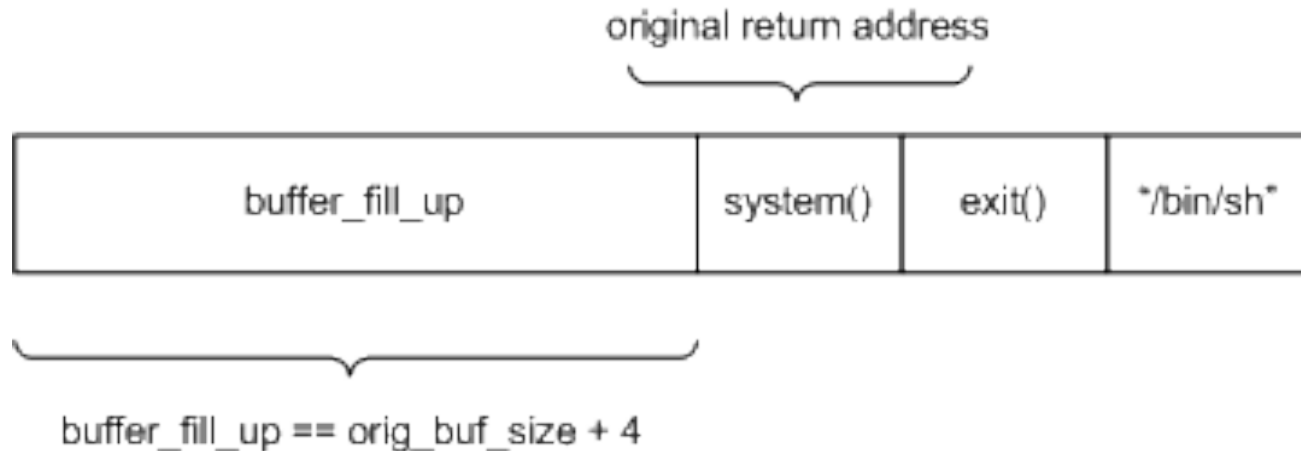


image stolen from: <http://www.w00z13.net/>

# The Exploit (DEP)

What if there is no library function available that will perform our task?

# The Exploit (DEP)

What if there is no library function available that will perform our task?

Do we even HAVE to return to functions?

# The Exploit (DEP)

Since we control the stack, we can "return" (jmp) to any location we desire.

# The Exploit (DEP)

Since we control the stack, we can "return" (jmp) to any location we desire.

What if we return to a "ret" instruction?

# The Exploit (DEP)

Now what if we return to a sequence of instructions ending in ret?

# The Exploit (DEP)

Now what if we return to a sequence of instructions ending in ret?

What if we do it again?



# The Exploit (DEP)

Now what if we return to a sequence of instructions ending in ret?

What if we do it again?

What if we continue indefinitely?

# The Exploit (DEP)

GADGETS!

# The Exploit (DEP)

## GADGETS!

- Think of `%rsp` as a makeshift instruction pointer
- intermingles code and data
- makes programming this "weird machine"\* interesting

\*shout out to pope

# The Exploit (DEP)

`int mprotect(void *addr, size_t len, int prot)`

- `void *addr`
  - The starting address of the page we wish to modify.
- `size_t len`
  - How many bytes are we modifying?
- `int prot`
  - What are we setting the protections to?

# The Exploit (DEP)

The calling convention is:

- argument 1: %rdi
- argument 2: %rsi
- argument 3: %rdx

...

# The Exploit (DEP)

Our ROP chain must:

- set %rdi to our shellcode's address
- set %rsi to (&shellcode - %rdi) + len(shellcode)
- set %rdx to +rwx (i.e., 7)
- return to &mprotect
- return/call/jmp to our shellcode

# The Exploit (shellcode)

- `gets()` is annoying to our file descriptor
- to get around this we'll use a bindshell

# The Exploit (shellcode)

```
sock = socket(AF_INET, SOCK_STREAM, 0);
bind(sock, &servaddr, sizeof(servaddr));
listen(sock, 1);
conn = accept(sock, NULL, NULL);
for (i = 0; i < 3; ++i)
    dup2(conn, i);
execve("/bin/sh", {"sh", "-i", NULL}, NULL);
```



# The Exploit (shellcode)

- Now that we know what we're doing, we can talk with the kernel directly instead of going through a library wrapper
- To do this, we need to lookup the syscall numbers

# The Exploit (shellcode)

```
$ vi arch/x86/syscalls/syscall_64.tbl
```

```
...
```

41	common	socket	sys_socket
49	common	bind	sys_bind
50	common	listen	sys_listen
43	common	accept	sys_accept
33	common	dup2	sys_dup2
59	64	execve	stub_execve

# The Exploit

demo

# Lessons

- Memory corruption is involved
- Modern mitigations require multiple bugs to gain code execution
- amd64 ABI doesn't require any pop-pop-rets but would be bad news if you lacked register control
- Manually writing ROP is either fun or the worst experience of your life

# Questions

Ask away!

# References

- <http://www.phrack.org/issues.html?issue=49&id=14>
- <http://www.phrack.org/issues.html?issue=58&id=4>
- <http://www.phrack.org/issues.html?issue=59&id=9>
- <http://cseweb.ucsd.edu/~hovav/papers/s07.html>
- <http://uninformed.org/index.cgi?v=9&a=4>
- <http://trailofbits.files.wordpress.com/2010/04/practical-rop.pdf>
- <http://www.suse.de/~krahmer/no-nx.pdf>
- <http://www.phrack.org/issues.html?issue=67&id=9>

# References

- <http://www.phrack.org/issues.html?issue=57&id=15>
- <http://www.phrack.org/issues.html?issue=57&id=5>
- <http://esec-lab.sogeti.com/post/2011/07/05/Linux-syscall-ABI>
- [http://msdn.microsoft.com/en-us/library/windows/hardware/ff561499\(v=vs.85\).aspx](http://msdn.microsoft.com/en-us/library/windows/hardware/ff561499(v=vs.85).aspx)

# Images (don't sue me)

- <http://wasntnate.com/2012/05/how-address-space-layout-randomization-aslr-was-beaten/>
- <http://www.w00zl3.net/>



# End

all files: <http://security.cs.rpi.edu/~candej2>

email: [napum@demigods.org](mailto:napum@demigods.org) (may or may not be down)

Notes:

- Hack the planet!
- dancetillyanosebleeds